A Fully Synthesizable, Low Voltage and Low Power Stochastic Flash A/D Converter with Wide Input Range

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Department of Information Engineering

The University of Kitakyushu

Author: Zou Xuncheng

(Advisor: Shigetoshi Nakatake)

CONTENTS

CONTENTS	I
1. Introduction	1
1.1 Background and motivation	
1.2 Thesis content and structure	10
2. Analog-to-Digital Converter	
2.1 Working Processes	12
2.1.1 Sample and Hold	13
2.1.2 Quantification	14
2.2 Main Specifications of ADC	15
2.2.1 Basic Specifications	15
2.2.2 Static Specifications	17
2.2.3 Dynamic Specifications	20
3. Stochastic Flash ADC (SFADC)	
3.1 Principle	24
3.2 Number of Comparators Required	27
3.2.1 Theoretical derivation	
3.2.2 Solution with uniformly distributed comparator offset	
3.2.3 Numerical Monte Carlo simulation	
3.3 Wallace Tree Adder	37

4. Fully Synthesizable Rail-to-rail Dynamic Comparator	
4.1 Review and Proposed Rail-to-rail Dynamic Comparator	41
4.1.1 Review	41
4.1.2 Proposed Rail-to-rail Dynamic Comparator	43
4.2 Simulation	47
4.3 Chapter Summary	55
5. Linearity Enhancement Technique (LET)	56
5.1 Prior Works	56
5.2 Our proposed Linearity Enhancement Technique (LET)	60
5.2.1 Description of LET	62
5.2.2 The reduction in the number of comparator groups	66
5.2.3 The determination of the number of comparators	
5.3 Effect of the proposed linearity enhancement technique	67
6. Implementation and Simulation	74
7. Conclusion	78
REFERENCES	

CHAPTER 1 Introduction

1.1 Background and motivation

Recent advances in technologies of wireless sensors. communication and embedded processors have enabled the design of small-size low-power and low cost devices that can be networked or connected to the Internet [1]. These are the key components of the emerging paradigm of Internet-of-things (IoT) [2,3]. A few examples of such applications are wireless sensor networks, biomedical and implantable devices/networks, ambient intelligence, wearable computing, smart grids, pollution monitoring, plant monitoring, smart warehouses [4-6]. The applications explicitly rely on the availability of sensor nodes that are energy autonomous and extremely small sized [7]. It can be said that the booming development of the Internet of Things is inseparable from the lowpower and miniaturization of electronic devices.

Low-power has emerged as a principal theme in today's electronics industry. The need for low power has caused a major paradigm shift where power dissipation has become as important a consideration as performance and area [8].

The power of a circuit is defined as (Eq. 1.1):

$$P(t) = i(t) \cdot v(t)$$
 (Eq. 1.1)

where i(t) is the instantaneous current provided by the power supply, and v(t) is the instantaneous supply voltage [9]. When we discuss the battery life or the energy dissipation of the system, we are more concerned about the average power consumption over a period of time. At this time, assuming that the voltage is constant, reducing the average current can reduce the average power consumption of the electronic device. In addition, intuitively, lowing the supply voltage can also reduce the power consumption of a circuit.

Specifically for CMOS circuits, power dissipation is caused by three sources: 1) the leakage current which is primarily determined by the fabrication technology, consists of reverse bias current in the parasitic diodes formed between source and drain diffusions and the bulk region in a MOS transistor as well as the subthreshold current that exists at the gate voltages below the threshold voltage, 2) the short-circuit current (crowbar current) which is due to the DC path between the supply rails during output transitions and 3) the charging and discharging of capacitive loads during logic changes[8].

The power consumption induced by the leakage current is also called static power consumption; The total power consumption caused by the short-circuit current and the charging and discharging of capacitive load is called dynamic power consumption. Static power consumption is the product of the device leakage current and the supply voltage. Total static power consumption, P_S , can be obtained as shown in (Eq. 1.2).

$$P_s = \sum (leakage \ current) \times (supply \ voltage)$$
 (Eq. 1.2)

The short-circuit power consumption component is less intuitive to be modeled because it depends on both the technology and the design parameters. It depends on the threshold and supply voltages, the drive strength of the gate, the frequency of operation, the input slope, and the output load connected to the gate [10]. A closed form for a symmetric inverter with the assumption of zero load capacitance at the output was derived was proposed as (Eq. 1.3) in [11]:

$$P_{sc} = \frac{\beta}{12} \cdot (V_{dd} - 2V_T)^3 \cdot \tau \cdot f_{clk} \quad \text{(Eq. 1.3)}$$

where P_{SC} represents the short-circuit power dissipation, β represents the strength of the transistors, V_T and V_{dd} are the threshold and supply voltages, respectively, τ is the input slope, and f_{clk} is the frequency of operation. It has also a strong dependency on the ratio between the supply and threshold voltages.

The short-circuit and leakage currents in CMOS circuits can be made small with proper circuit and device design techniques [8]. The dominant source of power dissipation is thus the charging and discharging of the node capacitances (also referred as the switching power dissipation) and is given by:

$$P_{sw} = 1/2 \cdot C \cdot V_{dd}^2 \cdot E(sw) \cdot f_{clk} \quad \text{(Eq. 1.4)}$$

Where C is the physical capacitance of the circuit, V_{dd} is the supply voltage, E(sw) (referred as the switching activity) is the average number of transitions in the circuit per $1/f_{clk}$ time, and f_{clk} is the clock frequency.

In summary, supply voltage scaling seems to be a good approach for power optimization, since the power normally yields considerable savings thanks to the strong dependence of power on supply voltage V_{dd} . However, the lower supply voltage means the lower circuit speed. Designers will have to make a trade-off between power consumption and circuit speed.

As the process continues to become more advanced, this contradiction has been alleviated. According to Moore's law [12], the number of transistors in an integrated circuit (IC) doubles about every two years, revealed in Fig. 1.1 [13]. The improvement of the process makes the feature size smaller. Also, the gate oxide of the MOSFET becomes thinner. The breakdown voltage of the device decreases, so the power supply voltage also decreases. The historic trend [14] in supply voltage is shown in Fig. 1.2.



Fig. 1.1: A semi-log plot of transistor counts for microprocessors against dates of introduction, nearly doubling every two years [13].



Fig. 1.2: Power-supply voltage as a function of feature size [14]

For digital circuits, the impact of low supply voltage on speed is offset by decreased gate capacitance. Therefore, process advance is conducive to power optimization for digital circuits. However, for analog circuits, the issue caused by low supply voltage can be big.

First of all, the primary index of the analog circuit is the signal-tonoise ratio. Lowering the supply voltage means that the signal swing is reduced, but the noise does not decrease in synchronization with the supply voltage. Therefore, low supply voltage has an adverse effect on the signal-to-noise ratio. Secondly, the threshold voltage of the device will not decrease synchronously with the supply voltage due to leakage, which makes the traditional circuit structure (such as cascode) no longer be adaptive at low supply voltage. Finally, since the threshold voltage does not decrease synchronously with the supply voltage, the operating region of the device is closer to the sub-threshold under low supply voltage, which is adverse to the linearity of the analog circuit.

In addition to the issues caused by the low supply voltage, the price of unit area in the advanced semiconductor manufacturing process is very high, but the overall size of analog circuits does not shrink proportionally with the reduction of feature size (especially passive devices, such as inductors, whose size has nothing to do with the feature size). Hence, the cost of analog circuits actually rises under the advanced semiconductor manufacturing process.

Consequently, digital circuits are more beneficial in scaling down than analog circuits. The use of digital circuits to replace analog circuits as much as possible can enjoy the benefits of process advances. Many related researches have been done, such as alldigital PLL, digital LDO, digital OPAMP, synthesizable RF transmitter [15-18].

Analog-to-digital converter (ADC) is an indispensable component in SOC. It almost represents the highest level in integrated circuit design, and has always been the focus and hot topic in IC field. Analog signals in the real world, such as temperature, pressure, sound and image, need to be converted to a digital form, which contributes to storing, processing or transmitting generally. ADC plays a role of a bridge between analog world and digital world. An analog signal comes into ADC, and a digital result would be generated.

ADC converts an analog value to a digital code according to specified rule. According to the difference in working mechanisms, it evolves into various of structures, including flash ADC, pipelined ADC, SAR ADC, $\Sigma\Delta$ ADC. Each has its own merits in speed, power, resolution, input bandwidth or other performance. Fig. 1.3 shows the comparison of several kinds of ADCs in regard to resolution and sampling rate. Compared to the other ADC, flash ADC is known for its high speed. Its resolution is usually between 4 to 9 bits, and the

operating frequency ranges from 10KHz to 10GHz. It is widely used in modulator, radio receiver, flash memory and so on.



Fig 1.3: Comparison of ADCs in regard to Resolution & Speed

However, it is difficult for conventional flash ADC to achieve higher resolution. The reason is that every time the resolution of one bit is increased, the number of comparators will double. Moreover, accomplishing the transformation from analog to digital needs comparator, which plays an important role in ADC. However, the device mismatch induced by process variation results in comparator offset, which affects the linearity of transfer function of ADC. As the process becomes finer, the feature size is increasingly reduced, making the offset voltage of the comparator more and more difficult to overcome, thus the comparator offset is becoming one of critical factors for designer to design a good ADC.

To deal with issues of the comparator offset, some of researches had been done to reduce or cancel the comparator offset using special techniques [19-21]. There are two kinds of methods in general. One is adopting large size device, which exchanges chip area for small offset. The other is utilizing calibration. However, both of them cost a large design overhead inevitably. Therefore, a stochastic ADC [22] comes into being, which aggressively makes use of the comparator offset variability rather than leaving nothing to do.

Targeting a flash architecture based on the idea of the stochastic ADC, the prior work in [23] has analyzed a cumulative distribution function (CDF) of offsets among a lot of comparators induced by the process variation, and presents a stochastic flash ADC (SFADC), proposing a conversion mechanism to employ an approximately linear section of the cumulative distribution function as the transfer function.

However, there are two main problems we have to deal with in the SFADC. First issue is power. In order to express the probability through the number of comparators whose output is logic one, the SFADC needs a mass of comparators (far more than a conventional flash ADC) to meet the statistic requirements. Hence, if single comparator is not power-effective, the overall power will be very

high. The second problem is input range. In the configuration of [23], input-output transfer function is approximately linear only within the input range from $-l\sigma$ to $+l\sigma$ of comparator offset standard deviation. The input range of the SFADC is limited due to the bad linearity. In order to improve the signal-to-noise ratio of the ADC, it is necessary to broaden the linear range as much as possible.

Therefore, this work focuses on SFADC, and revolves around how to solve these two problems. And it is committed to propose a low supply voltage, low power consumption SFADC with a wide input range.

1.2 Thesis content and structure

In view of the superiority of SFADC under the deep sub-micron process size, we present a fully synthesizable SFADC, which can operate at the supply voltage of 0.6V with power consumption as low as 1.5mW at the clock frequency of 250MHz. By employing the all-digital comparator, the SFADC can be described with Verilog language and synthesized according to a standard digital design flow. Cross-coupled dynamic comparator structure saves the overall power due to remarkable control of dynamic power consumption. In addition, the rail-to-rail characteristic of comparator and the proposed linearity enhancement technique based on SFADC, allow us to design a wide-range SFADC. The rest of the thesis is organized as follows. Chapter 2 introduces ADC's theoretical overview and some of significant specifications. Chapter 3 mainly introduces the fundamental and properties of SFADC, and discusses the relationship between the number of comparators and signal-to-noise-and-distortion ratio. Chapter 4 is devoted into describing the fully synthesizable rail-to-rail dynamic comparator in our work. Chapter 5 elaborates the proposed linearity enhancement technique (LET). Chapter 6 illustrates the structure of the proposed SFADC and shows the advantage on power consumption, supply voltage and input-range based on simulation results. Chapter 7 concludes this work.

CHAPTER 2 Analog-to-Digital Converter

2.1 Working Processes

From an analog value which is continuous in time and amplitude, to a digital value which is discrete in time and amplitude, signal usually needs to go through four processes -- pre-filtering, sampling, quantization, and encoding, as shown in Fig. 2.1. Pre-filter filters the parts of signal outside the Nyquist frequency, to avoid aliasing caused by high frequency signal in the baseband of AD converter. Hence the pre-filter also can be called 'anti-aliasing filter'. A S&H (Sample and Hold) circuit is connected to the anti-aliasing filter. A sampling circuit produces a sequence of δ functions of which amplitudes are equal to the ones of signal at the sampling times. Next, the holding circuit maintains the sampling signal, and makes it remain unchanged during the transformation. Then quantizer divides reference voltage into 2^{N} -1 sub-domains (N is the number of digital output bits). After finding out which sub-domain the sampling signal corresponds to, digital encoder begins to encode and outputs a digital result. In a transformation period, a sampling of the analog input signal is converted to an equivalent digital output code.



Fig. 2.1: A basic working process of ADC

2.1.1 Sample and Hold

Sampling is an operation of converting continuous time signal to discrete time signal. Ideally, sampling operation generates a pulse sequence whose amplitudes are equal to signal's amplitudes at sampling points. Then holding operation maintains the sampling signal, and makes it remain unchanged during the transformation, generating an analog signal $x_s(n^*T_s)$ whose time is discrete and amplitude is continuous, as shown in Fig. 2.2. T_s is the sampling cycle and n^*T_s (n = 0, 1, 2, ...) is called sampling moment.



Fig. 2.2: Sample and hold

Assume that f_m represents highest frequency of input signal, and $1/(2*f_m)$ is called Nyquist interval. Only if the sampling interval T_s is less than the Nyquist interval $1/(2*f_m)$, the sampling signal $x_s(n*T_s)$ can be restored to the original analog input signal x(t), which is called Nyquist–Shannon sampling theorem.

2.1.2 Quantification

Quantification is an operation of converting continuous amplitude signal to discrete amplitude signal. Taking an ideal 3-bit ADC as an example, its transfer function is shown in Fig. 2.3.



Fig. 2.3: Transfer function of a 3-bit ADC in quantification

Input signal is an analog value from 0 to the *FS* (Full Scale), and the output is a group of quantitative levels, 8 (2³) in total, expressed in the digital discrete form. A *N*-bit ADC has 2^N digital codes, so *FS* is divided into 2^N subintervals. The width of subinterval is defined as the least significant bit (LSB), namely $1LSB = FS/2^N$. The trip points of transfer function are at x = k * LSB (k = 1, 2, ...). Similarly, digital to analog conversion also can be easily implemented. Each digital bit (b_x) has a weight of 2^{x-1} , and a digital value can be restored into an analog value as (Eq. 2.1).

$$V_{analog} = LSB^{*}(b_{N}^{*}2^{N-1} + b_{N-1}^{*}2^{N-2} + \dots + b_{2}^{*}2^{1} + b_{1}^{*}2^{0}) + LSB/2$$

(Eq. 2.1)

2.2 Main Specifications of ADC

2.2.1 Basic Specifications

(1) Sampling Frequency

Sampling frequency is the reciprocal of the sampling time, and represents the times of conversion from continuous analog signal to discrete digital signal per second in the AD converter. The unit is Hz.

(2) Resolution

Resolution is the number of digital bits to represent an analog input. The resolution, together with reference voltage determines the minimum detectable voltage.

(3) Nyquist Frequency

Nyquist frequency, decided by Nyquist–Shannon sampling theorem, is the maximum bandwidth of input signal, which is equal to half the sampling frequency. If the maximum frequency of input signal is beyond the Nyquist Frequency, energy information of signal at different frequency can't be restored with digital output signal.

(4) Input Signal Range

Input signal range is the quantitative range of ADC, generally determined by reference voltage. If the input signal is beyond the range of input, the AD converter will cause distortion.

(5) Power Consumption

Power consumption is energy consumed by AD converter per unit time when ADC works. Now low power consumption has become an important index in the ADC designing.

(6) Area

is the area of ADC on chip, often expressed in mm², which determines the cost of ADC.

2.2.2 Static Specifications

(1) Offset

The offset describes an output shift for zero input, also expressed as shift error. It is the skewing of transfer function of ADC, expressed in mV or percentage of full scale. As shown in Fig.2.4, all of quantization levels shift an offset error.



Fig. 2.4: Offset error

(2) Gain Error

Gain error is the difference between the ideal analog input signal and the actual analog input signal when input causes a transition to full scale, expressed in the mV or the percentage of full scale, as shown in Fig. 2.5. Another measure of gain error is the error on slope of transfer characteristic curve relatively to the ideal characteristic curve around the origin of coordinates. Unlike DNL and INL, both of gain error and offset error are linear error.



Fig. 2.5 Gain error

(3) Quantization Error

Quantization error is defined as the difference between ideal *N*-bit ADC output and infinite resolution converter's output, also known as the least effective bit error. The analysis of quantization noise seems difficult since the quantization noise is a function of the input signal. Fortunately, under some specific conditions, quantization error can be approximate to white noise irrelevant to the input signal. For an ideal AD converter, the input and output are not one to one correspondence. When the input changes, the output may not change. Quantization error is innate in the AD converter, and is the lower limit of ADC error. Usually the only way to reduce the quantization error is improving resolution.

(4) Common Mode Error

is applied to differential inputs ADC, which represents the changes on output code when common mode input changes by a given value. It is usually measured by measuring the changes on equal common mode inputs when output code changes by 1 LSB, expressed in LSB generally.

(5) Differential Non-linearity Error (DNL)

is the step difference between actual transfer function and ideal one, expressed usually in LSB. Often the maximum DNL is simply referred as DNL. Differential non-linearity error can be expressed by (Eq. 2.2).

$$DNL_x = \frac{V_{LSB,x} - V_{LSB}}{V_{LSB}} \quad (Eq. 2.2)$$

where $V_{LSB,x}$ is the actual LSB

(6) Integral Non-linearity Error (INL)

Integral non-linearity error (INL) is the deviation of actual transfer function from endpoint fit line, usually expressed in LSB. Often the maximum INL is simply referred as INL. Integral non-linearity error can be expressed by (Eq. 2.3).

$$INL_x = \sum_{i=0}^{x} DNL(i) \quad (Eq. 2.3)$$

2.2.3 Dynamic Specifications

(1)Input Impedance

is the impedance between ADC inputs. The input impedance performs as resistance at low frequency. Ideally, input impedance of ADC is infinite when inputting voltage, and zero when inputting current. At high frequency the input impedance is usually determined by capacitive devices. As the switched capacitors are usually used in the ADC sampling, input impedance of ADC should match to input terminals at high frequency.

(2) Signal-to-noise Ratio (SNR)

Under the specific input and sampling frequency, ratio between the ADC output signal power and noise power is defined as signal-tonoise ratio. It is the ratio of ADC output signal and noise without consideration of the distortion. For an ideal ADC, the SNR can be expressed as (Eq. 2.4).

$$SNR = 10 \log \left(\frac{Signal}{Noise}\right) (db)$$
 (Eq. 2.4)

The input signals are generally sine waves when measuring. When only the quantization noise is considered, the signal-to-noise ratio of an ideal *N*-bit ADC can be calculated as (Eq.2.5).

$$SNR = 6.02N + 1.76 (db)$$
 (Eq. 2.5)

(3) Signal-to-noise-and-distortion Ratio (SNDR)

Signal-to-noise-and-distortion ratio is referred to the ratio between the AD converter output signal power and the sum of all noise and harmonic power, generally expressed in dB. The calculation formula is revealed as (Eq. 2.6).

$$SNDR = 10 \log \left(\frac{Signal}{Noise + Harmonics} \right) (db)$$
 (Eq. 2.6)

(4) Total Harmonic Distortion (THD)

Total harmonic distortion is the ratio between total harmonic power and fundamental wave power within a specific frequency range, defined as (Eq. 2.7).

$$THD = 10 \log \left(\frac{Total \ Harmonic \ Distortion}{Signal} \right) (db)$$

(Eq. 2.7)

(5) Effective Number of Bits (ENOB)

ENOB of AD converter is a dynamic value changing with signal frequency, and it reflects the effective conversion bits at different signal frequency in dynamic working.

With the increasing of input frequency, the overall noise (also distortion) will increases, thus reducing the ENOB and SNDR. So ENOB is often defined by SNDR as (Eq. 2.8).

$$ENOB = \frac{SNDR-1.76}{6.02}$$
 (Eq. 2.8)

and SNDR is expressed in dB.

(6) Dynamic Range

is the value of input signal when SNR (or SNDR) becomes zero, usually expressed in dBFS (Full Scale).

(7) Spurious Free Dynamic Range (SFDR)

SFDR is defined as the ratio between RMS (root mean square) of input signal amplitude and RMS of the largest distortion component in first Nyquist domain. It is the difference value between component of fundamental wave (expressed in dB) and the highest noise component (expressed in dB) in the output spectrum. SFDR depends on the amplitude of input signal. For large input, the highest distortion is usually a harmonic component, but when the signal amplitude gets smaller, the distortion caused by the signal can be ignored. At this time, the distortion is often determined by other source rather than input signal. It is expressed as (Eq. 2.9).

$$SFDR = 10 \log \left(\frac{Signal}{Largest Spurious or Harmonic}\right)$$
 (db) (Eq. 2.9)

For any good INL AD converter, SFDR is greater than SNR. Noise and harmonic are main factors limiting dynamic range of AD converter, thus SFDR is a very important specification for AD converter.

(8) Effective Resolution Bandwidth (ERBW)

Effective resolution bandwidth is referred to the frequency of input signal when SNDR of ADC decreases by 3 dB relatively to low frequency.

(9) Figure of Merit

Above specifications of ADC affect each other, so it is difficult to use a particular specification to measure the overall performance of the ADC. Now a common practice is taking the main performance parameters of ADC as a comprehensive index, and it is called FoM (Figure of Merit). The smaller FoM value is, the better the performance is. It can be calculated as (Eq. 2.10).

$$FoM = \frac{Power}{2^{ENOB} * f_s} \quad (Eq.2.10)$$

Power is the total Power consumption of ADC; f_s is the sampling frequency of ADC.

CHAPTER 3 Stochastic Flash ADC (SFADC)

3.1 Principle

In a conventional flash ADC, input signal is connected to input ports of a group of comparators. The reference voltage of each comparator is set precisely by reference ladder, so that all comparator thresholds are equally spaced by 1 LSB.

However, different from conventional one, as shown in Fig. 3.1, in the SFADC, an input signal line and a common reference voltage are connected to the inputs of comparators. Plus, following on the heels of a group of comparators, a Wallace Tree adder [24] is used to sum up the number of logic '1' from all comparator outputs.



Fig 3.1: Wallace Tree Adder

In a SFADC, the occurrence of comparator offset is not a drawback any longer, but rather something available. Due to device mismatch or processing variation, comparator offset appears to be random. The variation of comparator offset can be assumed to be a Gaussian distribution with a mean (μ) of zero and variance (σ^2) which is inversely proportional to the comparator area [25], as shown in Fig. 3.2(a). When a reference voltage is applied to the comparator, the comparator offset's probability density function (PDF) becomes the PDF of Gaussian distribution with a mean of reference voltage, expressed by (Eq. 3.1).

$$f(x;\mu,\sigma) = \frac{1}{\sigma\sqrt{2\pi}} \exp\left[-\frac{(x-\mu)^2}{2\sigma^2}\right] \quad (\text{Eq. 3.1})$$

The Wallace Tree adder counts for the number of high level voltage (i.e. logic '1') in the outputs of comparators. As a result, when a ramp signal is given to a SFADC, the output of Wallace Tree adder is the cumulative distribution of a Gaussian distribution with a mean of reference voltage, as shown in Fig.3.2(b). In other words, the transfer function of a SFADC is regarded to be the cumulative distribution function (CDF) of the random comparator offsets, expressed as (Eq. 3.2).

$$F(x,\mu,\sigma) = \frac{1}{2} \left[1 + \operatorname{erf}\left(\frac{x-\mu}{\sigma\sqrt{2}}\right) \right] \quad (\text{Eq. 3.2})$$

where

$$\operatorname{erf}(x) = \frac{2}{\sqrt{\pi}} \int_0^x e^{-t^2} dt$$

If the amount of comparators is large enough, the PDF of comparators offset can be approximated to a Gaussian distribution function, as well as the conversion function is the CDF of Gaussian distribution. The mid-point of the conversion function is precisely corresponding to the reference voltage. In general, a conventional *N*-bit flash ADC needs $2^{N}-1$ comparators, however, far more comparators must be incorporated into a SFADC to get closer to CDF of Gaussian distribution. SFADC's transfer function has a good linearity between -1σ to $+1\sigma$ (σ is the standard deviation of comparator offset). Thus the SFADC without calibration generally works in this input range.



Fig. 3.2 (a): Comparator offset's distribution (b)Transfer function of SFADC

3.2 Number of Comparators Required

3.2.1 Theoretical derivation

Assuming that for any comparator *i*, the comparator offset's probability density function is $P_{offset,i}$. When an input voltage of *v* is given to it, the output D_i follows the Bernoulli distribution, whose probability mass function $P_v(Di)$ is expressed as (Eq. 3.3).

$$P_{\nu}(D_i) = F_i(\nu)^{D_i} \cdot [1 - F_i(\nu)]^{1 - D_i} \quad (\text{Eq. 3.3})$$

where

$$F_i(v) \equiv \int_{-\infty}^{v} P_{offset,i}(V) dV$$

and $D_i = \{0, 1\}$.

If a SFADC contains N comparators in total, since the offsets of each comparator are independent of each other, the sum of the Ncomparator outputs (that is, the total output code D) represents the sum of N independent Bernoulli trials. D is given by (Eq. 3.4),

$$D = \sum_{i=1}^{N} D_i$$
 (Eq. 3.4)

which follows a Poisson binomial distribution.

The exact PDF of the total output code D can be calculated by convolving the output codes of all the comparators, which is a huge

computation of the order $O(2^N)$ [26]. Even so, we can still use the first 2 moments to approximate the PDF of *D* to a curve in the Pearson family of distributions [27]–[28]. The first and second moment about 0 of *D* is given as (Eq. 3.5) and (Eq. 3.6), respectively.

$$M_{1} = \sum_{D=0}^{N} P_{\nu}(D) \cdot D = \mu_{D} = \sum_{i=1}^{N} F_{i}(\nu) \quad (\text{Eq. 3.5})$$
$$M_{2} = \sum_{D=0}^{N} P_{\nu}(D) \cdot D^{2} = \sum_{i=1}^{N} F_{i}(\nu) \cdot [1 - F_{i}(\nu)] + [\sum_{i=1}^{N} F_{i}(\nu)]^{2} \quad (\text{Eq. 3.6})$$

where M_1 also represents the mean μ_D of output code D at the input voltage of v.

Only when all the comparators have the same P_{offset} , the distribution of the total output code D will become a binomial distribution. At this time, the output code D has a probability mass function of (Eq. 3.7).

$$P_{\nu}(D) = \binom{N}{D} F(\nu)^{D} [1 - F(\nu)]^{N-D} \quad (\text{Eq. 3.7})$$

where

$$F(v) \equiv \int_{-\infty}^{v} P_{offset}(V) dV$$

Then the first and second moment about zero of D is simplified to (Eq. 3.8) and (Eq. 3.9), respectively.

$$M_1 = \mu_D = N \cdot F(v)$$
 (Eq. 3.8)

$$M_2 = N \cdot F(v) \cdot [(N-1) \cdot F(v) + 1]$$
 (Eq. 3.9)

where μ_D is the average output code, which is a function of the input *v*. In fact, it represents the average transfer function of SFADC.

It should be emphasized that (Eq. 3.5) and (Eq. 3.6) are the general case of (Eq. 3.8) and (Eq. 3.9), regardless of whether the PDF of the comparator offset is same.

For an SFADC without calibration, only the relatively linear region in the transfer function can be used for ADC conversion. Let L(v) be the ideal linear portion in μ_D , that is, the expected output.

Next we define the quantization error q_e as the difference between the expected output L(v) and the actual output D, namely

$$q_e = L(v) - D \quad \text{(Eq.3.10)}$$

Therefore, the variance of the quantization error q_e can be derived as follows.

$$E[q_e^2] = \sum_{D=0}^{N} P_v(D) \cdot (L(v) - D)^2$$

= $\sum_{D=0}^{N} P_v(D) \cdot (L(v)^2 - 2 \cdot D \cdot L(v) + D^2)$
= $L(v)^2 \cdot \sum_{D=0}^{N} P_v(D) - 2 \cdot L(v) \cdot \sum_{D=0}^{N} D \cdot P_v(D) + \sum_{D=0}^{N} D^2 \cdot P_v(D)^{\text{(Eq. 3.11)}}$
= $L(v)^2 - 2 \cdot L(v) \cdot M_1 + M_2$

Next, let's discuss three special cases.

(1) When all comparators have the same PDF of F(v), (Eq. 3.11) can be reduce to (Eq. 3.12) combining (Eq. 3.8) and (Eq. 3.9).

$$E[q_e^2] = (N^2 - N) \cdot F(v)^2 - [2 \cdot N \cdot L(v) - N] \cdot F(v) + L(v)^2$$
(Eq. 3.12)

(2) When only the linear portion about v is contained in μ_D , in other words, when L(v) is equal to μ_D , (Eq. 3.11) can be simplified to (Eq. 3.13) combining (Eq. 3.5) and (Eq. 3.6).

$$E[q_e^2] = \sum_{i=1}^{N} F_i(v) \cdot [1 - F_i(v)] \quad (\text{Eq. 3.13})$$

(3) When the conditions of (1) and (2) are met at the same time,(Eq. 3.11) can be further reduced to (Eq. 3.14).

$$E[q_e^2] = N \cdot F(v) \cdot [1 - F(v)]$$
 (Eq. 3.14)

The quantization noise energy Q is expressed as the integral of $E(q_e^2)$ over the entire input range R, namely

$$Q = \int_{R} E[q_e^2] dv \quad \text{(Eq. 3.15)}$$

A ramp signal can be described as a random variable with a uniform PDF [29]. The variance of a random variable is equivalent to its mean-square power, and since the variance of a uniform PDF is found to be

$$Var(PDF_{uniform}) = \frac{\Delta^2}{12}$$
 (Eq. 3.16)

where Δ is the range of the PDF.

The signal energy S in the output is uniformly distributed between L(a) and L(b), where a and b are the start and end points of the linear input range R respectively. The signal energy S in the output is equal to the integral of the mean-square power over the entire input range of [a, b].

$$S = \frac{[L(b) - L(a)]^2}{12} \cdot (b - a) \quad \text{(Eq. 3.17)}$$

Finally we can calculate the SNDR as (Eq. 3.18) combining (Eq.3.15) and (Eq. 3.17).

$$SNDR = \frac{S}{Q} = \frac{(b-a) \cdot [L(b) - L(a)]^2}{12 \int_a^b E[q_e^2] dv}$$
(Eq. 3.18)

Although (Eq. 3.18) is still very complicated, we can draw some useful conclusions. Firstly, for both L(v) and $E[q_e^2]$, according to the definition of L(v) and (Eq. 3.11), it is not difficult to find that they are all related to μ_D , which is uniquely determined by the number of comparators N and the offset's distributions $P_{offset,i}$. Secondly, both numerator and denominator of (Eq. 3.18) are related to the input range [a, b]. Therefore, SNDR is decided by the input signal range, the distribution of comparator offset, and the number of comparators.

In addition, since SFADC usually works in a linear input range, for a given linear input range [a, b], L(v) can be replaced by approximately μ_D . Thus the numerator *S* is proportional to N^2 . In the same way, $E[q_e^2]$ can be calculated by (Eq. 3.13), so the denominator *Q* is proportional to *N*. As a result, the ratio of *S* and *Q*, SNDR is proportional to the number of comparators *N*. It indicates that every time the number of comparators is increased to 4 times of the original number, SNDR is also increased to 4 times. When converted into dB, SNDR_{dB} will increase by 6dB. According to (Eq.2.8), the ENOB is increased by one bit.

3.2.2 Solution with uniformly distributed comparator offset

The general solution from (Eq. 3.18) is applied to two special cases in which comparator offset is uniformly distributed.

(1) Assume that comparator offset is randomly and uniformly distributed along full-scale from 0 to 1, namely

$$P_{offset}(v) = 1/\Delta = 1$$
 (Eq. 3.19)
 $F(v) = v$ (Eq. 3.20)

Then according to (Eq. 3.7), the probability mass function of the output code is

$$P_{\nu}(D) = \binom{N}{D} (\nu)^{D} \cdot (1-\nu)^{N-D} \quad \text{(Eq. 3.21)}$$

When a random uniformly distributed input is applied to this SFADC, according to (Eq. 3.8), the mean of the output code is N^*v ,

that is, the average transfer function is $\mu_D = N^* v$. It can be seen that the average transfer function does not have a non-linear portion about v, so the expected output code is also $L(v) = N^* v$.

The variance $E[q_e^2]$ of the quantization error q_e is found in (Eq.3.14), that is, $N^*v^*(1-v)$. According to (Eq. 3.15) and (Eq.3.17), the signal power and quantization noise power contained in the output are $N^2/12$ and N/6, respectively. So SNDR is calculated by SNDR=N/2. It is consistent with the conclusion in [29].

(2) The comparators are equally divided into two groups, and their offsets are uniformly distributed on the input range [0,1/2] and [1/2,1], respectively.

Assuming that the PDFs of the offset in the two groups are $P_{offset1}$, $P_{offset2}$, respectively.

$$P_{offset1}(v) = \begin{cases} 2, & 0 < v < 1/2 \\ 0, & Other \end{cases}$$

$$P_{offset2}(v) = \begin{cases} 2, & 1/2 < v < 1 \\ 0, & Other \end{cases}$$
(Eq. 3.22)
(Eq. 3.23)

The CDFs of the offset in the two groups can be also calculated as follows.

$$F_{1}(v) = \int_{-\infty}^{v} P_{offset1}(V) dV = \begin{cases} 0, & v \le 0\\ 2v, & 0 < v \le 1/2\\ 1, & 1/2 < v \end{cases}$$
(Eq. 3.24)
$$F_2(v) = \int_{-\infty}^{v} P_{offset 2}(V) dV = \begin{cases} 0, & v \le 1/2\\ 2v - 1, & 1/2 < v \le 1\\ 1, & 1 < v \end{cases}$$
(Eq.3.25)

According to (Eq. 3.5), the mean of the total output code can be obtained as (Eq. 3.26).

$$\mu_D = \begin{cases} 0, & v \le 0\\ N \cdot v, & 0 < v \le 1\\ N, & 1 < v \end{cases}$$
(Eq. 3.26)

It can be seen that μ_D is exactly the same as the situation in case(1).

Also, the variance of the quantization error is derived from (Eq. 3.13).

$$E[q_e^2] = \begin{cases} N \cdot v \cdot (1 - 2v), & 0 < v \le 1/2 \\ N \cdot (2v - 1) \cdot (1 - v), & 1/2 < v \le 1 \\ 0, & Other \end{cases}$$
(Eq. 3.27)

According to (Eq. 3.15) and (Eq.3.17), the signal power and quantization noise power contained in the output are $N^2/12$ and N/12, respectively. The quantization noise power in case (2) is only half of case (1), though they have the same average transfer function μ_D . It reveals that, though having the same μ_D can determine that they have the same signal power, it cannot determine the variance of the quantization noise. Dividing the comparators with the same offset distribution into comparator groups with different distributions results in a change in the variance of the quantization noise of the SFADC, thereby affecting the SNDR.

3.2.3 Numerical Monte Carlo simulation

Although many prior works [23,25,29-31] have tried to establish a mathematical model of the relationship between SNDR and the number of comparators, the actual SNDR obtained in some works such as [29], does not match the SNDR calculated through the mathematical model. Since on the occasion of [29], multi-group structure of SFADC is adopted, but the effect of grouping to the quantization noise power is omitted.

Nevertheless, when the number of comparator groups is large, using (Eq. 3.13) to calculate the quantization noise power, so as to obtain the relationship between SNDR and the required number of comparators is still complicated.

Therefore, we use Monte Carlo simulation instead of mathematical model in the case of a large number of groups to establish the relationship between SNDR and the number of comparators required. We first take samples of random variable with a given distribution, and using these values as the trip points for an ideal SFADC, and after applying a full-scale sine input, the SNDR can be calculated through 4096-point FFT. Repeating this test 100 times allows us to find the average SNDR for a given number of comparators. Fig. 3.3 reveals the relationship between SNDR and the number of bits *N*, where the number of comparators is equal to 2^{N} -1. It can be seen that, to achieve an improvement of 6dB on SNDR, the number of comparators needs to be quadrupled, which is consistent with the

conclusion in the Section 3.2.1, also the conclusions of [23,25,29-31]. Although their theoretical value of the number of comparators is different, the conclusion that ENOB is proportional to 4^N is consistent.



Fig. 3.3: Averaged SNDR versus number of bits of SFADC

3.3 Wallace Tree Adder

The encoder counts the number of high levels output by the comparator bank and converts the result into a binary code form. The fault tolerance and speed must also be considered when designing the encoder. Working for summation of logic '1', Wallace tree adder is applied to encoding of SFADC. The encoder is realized by cascading the full adder FA into a Wallace Tree structure. Fig.3.4 shows an example of 6-bits Wallace Tree encoder [32]. Assuming that the number of full adders required for an *N*-bits encoder is X_N . Two *N*-bits encoders and *N* full adders forms a new (N+1)-bit encoder.

Hence,

$$\begin{cases} X_{N+1} = 2X_N + N \text{ (for } N \ge 2) \\ X_2 = 1 \end{cases}$$
 (Eq. 3.28)

By solving the recursive sequence (Eq. 3.28), the number of full adders required for an *N*-bits encoder is given by

$$X_{N} = \sum_{i=1}^{N} (i-1) \cdot 2^{N-i} \quad \text{(Eq. 3.29)}$$

With the properties of full adder, carry input on the full adder can be equated to an addend. A full adder can therefore be regarded as a three-input adder. In accordance with weight, bits of the same weight can be connected to the common adder in the next level for a summation while the bits of different weights are separated.



Fig. 3.4: An example of 6-bit Wallace tree adder

There are three advantages with this structure. Firstly, tree topology can reduce the transmission distance of digital signal, reducing the parasitic capacitance. Secondly, the structure of the Wallace tree adder doesn't concern on the order of the comparators. So Wallace tree adder structure can be flexibly designed, which makes it especially suitable for SFADC. Thirdly, the tree topology can be easily implemented in pipelined work. Inserting D-triggers to the critical path of the circuit, allows the Wallace tree adder to work at a high speed, shown as Fig. 3.5.

The maximum sample frequency is decided by adder cell delay when all of FAs are pipelined:

$$f_{\rm max} = 1/(adder \ cell \ delay)$$
 (Eq. 3.30)



Fig. 3.5: Pipelined work in Wallace tree adder

CHAPTER 4

Fully Synthesizable Rail-to-rail Dynamic Comparator

Due to large-scale adoption of comparators, decreasing the power consumption of single comparators becomes an effective method to decrease the overall power consumption of SFADC. In addition, digital circuits are more profitable from scaling down compared to analog circuits. They can obtain low parasitic capacitance, supply voltage and power consumption effortlessly in this context. The synthesizable characteristic saves them from the effort for customized layouts as well. Therefore, we focus on the realization of a fully synthesizable dynamic comparator with a wide input range.

The all-digital design of comparator have been concerned. [25] proposes a dynamic comparator based on two cross-coupled 3-input NAND gates. However, the common-mode input range(CMR) is limited to a high voltage region to avoid PMOS current more robust than the NMOS current of input transistors. To address the issue of the narrow CMR, two types of methods have been proposed in previous works.

[33]–[35] introduce additional pull-down networks. When the common-mode input is a low voltage, the strong sinking current flowing via the pull-down network offsets the effect of the PMOS

current, so that low input voltages are applicable. Nevertheless, the strength of the pull-down network needs careful consideration. The strength which is too strong or too weak cannot force output to the supply rail properly. Also, the method of utilizing a pull-down network causes high power consumption due to a severe shortcircuit current, which is disadvantageous for low-power design.

[36] and [37] combine the NAND-based comparator and the NORbased comparator, utilizing the complementary features of them. When the common-mode input is close to *VDD* or *VSS*, there is always a comparator that works correctly. Through a selection mechanism, the comparator with the valid result transmits data to the final comparator output. However, the alternative mechanism may lead to incorrect output results when the CMRs of the two comparators do not intersect and the input is just falling in the gap. Besides, the robust short-circuit in their decision phase results in low power efficiency.

4.1 Review and Proposed Rail-to-rail Dynamic Comparator

4.1.1 Review

[25] firstly proposes a dynamic comparator composed of standard cells. 3-input NAND gate is usual in the standard cell library. The manner of constructing the comparator with two NAND gates is

illustrated as Fig. 4.1. Transistors *MP5*, *MP6*, *MN5* and *MN6* constitute a regenerative latch. When the clock is low(reset phase), *MP3* and *MP4* are turned on; *MN3* and *MN4* are turned off. Thus output nodes *Vout+* and *Vout-* are precharged to *VDD*. When the clock is switched to high(decision phase), *MP3* and *MP4* are turned off, thus disabling the precharge of *Vout+* and *Vout-*. Due to the drain current difference between *MN1* and *MN2* induced by input voltage difference, *Vout+* and *Vout-* drop voltage at a different speed. As *Vout+* and *Vout-* are gradually pulled down, gate voltage on the one of PMOSs in the regenerative latch is first to reach *VDD-Vth*, thus enabling this PMOS to charge to corresponding output nodes. Then positive feedback effect starts and amplifies the voltage difference of output nodes. Finally, the output nodes are forced to the supply rails.



Fig. 4.1: NAND-based Comparator. (a) Schematic; (b)Symbol

However, a problem arises when common-mode input is close to VSS. In this case, *MP1* and *MP2* are turned on, producing a stronger current than *MN1* and *MN2*, which causes charging to output node continuously. Finally, both of the outputs are pulled up to high voltage, comparator failing to compare the two inputs. This phenomenon dramatically limits the CMR of the comparator.

4.1.2 Proposed Rail-to-rail Dynamic Comparator

In order to broaden the input range of the NAND-based comparator, we propose a block-based method to construct a dynamic comparator that is made up of two OAI211 gates(OR-AND-INVERTER). This cell is also typical in the standard cell libraries of advanced processes. Connect two OAI211 gates to construct our proposed comparator, as shown in Fig. 4.2. Different from the NAND-based comparator, *MP3*, *MP4*, *MN3* and *MN4* are "additional transistors" (marked in red).



Fig. 4.2: OAI211-based Rail-to-rail Dynamic Voltage Comparator (ORDVC). (a) Schematic; (b)Symbol

MP3 and *MP4* play the roles of valves controlled by the output voltage, cutting off the direct connection between the supply rail and the output nodes. The one whose gate voltage first reaches *VDD* in the decision phase completely blocks the current flowing from supply power to the corresponding drain. Therefore, *MP3* and *MP4* effectively avoid both output nodes to be pulled up to high voltages together when common-mode input is low. Since dynamic current only flows during regeneration, this structure is power-effective.

Taking advantage of the complementary idea, we further propose a variant composed of OAI211-based comparator and AOI211-based (AND-OR-INVERTER) comparator for a wider CMR, illustrated as Fig. 4.3. Somewhat different from the single comparator like OAI211-based comparator, two complementary comparators are

cross-coupled with each other, that is, they connect the gates of the additional transistors to each other's output nodes. The complementary signal *NCLK* of the clock signal *CLK* is obtained through an inverter. When the common-mode input voltage is close to *VDD*, the OAI211-based portion dominates; otherwise, the AOI211-based portion dominates when the common-mode input voltage is close to *VSS*.



Fig. 4.3: Merged Rail-to-rail Dynamic Voltage Comparator (MRDVC). (a) Schematic; (b)Symbol

Two variants have different initial conditions when they are switched to the decision phase, which predictably brings some differences in performances. For convenience, the OAI211-based rail-to-rail dynamic voltage comparator is called 'ORDVC' for short, and the merged version consist of the two complementary comparators is 'MRDVC'.

In terms of ORDVC in Fig. 4.2, *MP3* and *MP4* are cut off while *MN3* and *MN4* are turned on initially. On the one hand, the pulldown network is strengthened, and the pull-up network is weakened compared to a NAND-based comparator. The voltage of output nodes can swiftly meet the triggering condition of the positive feedback, which brings improvement on propagation delay. On the other hand, due to high voltage on the gates of *MN3* and *MN4* initially, device mismatch between *MN3* and *MN4* can lead to a large mismatch current. It takes a tremendous effort for the input to offset such effect, which makes comparators have a larger variation in comparator offset.

As far as MRDVC is concerned in Fig. 4.3, at the beginning of decision phase, *MP1*, *MP2*, *MN3* and *MN4* are turned on; on the contrary, *MN1*, *MN2*, *MP3* and *MP4* are cut off. At the moment, the merged comparator is shaped like an individual NAND-based comparator and NOR-based comparator. Therefore, it offers a legitimate mix of propagation delay and comparator offset variation.

4.2 Simulation

Based on the 65nm CMOS process, the prior art (Seo [33] and Aiello [36]) are replicated to compare with our proposed comparators through the simulation. Seo [33] is the most power-effective representative of the pull-down method, which adopts a power down logic for energy saving, while Aiello [36] is the representative of the complementary method. To facilitate the comparison, we use the minimum size of transistors and appropriate latches in all candidates. It is noted that for MRDVC, either *Vo1*, *Vo2*, or *Vo3*, *Vo4* can be used as complementary outputs connected to an appropriate latch, with a little influence on the simulation result.

Under the different supply voltages, the relationship between the propagation delay and common-mode input voltage is shown in Fig.4.4 when the load is 5fF and the differential voltage is 5mV. In regard to CMR, MRDVC is the only one covering a full rail-to-rail CMR under three different supply voltages. Aiello [36] fails to output the correct results at 120 mV when the power supply is 300mV, leading to a discontinuous CMR. Seo [33] only allows the common-mode voltage higher than 1/2 *VDD*. ORDVC only accepts about half of the rail-to-rail CMR in the cases of *VDD* equal to 300mV and 600mV.



Fig. 4.4: Propagation delay versus common-mode voltage under different supply voltages. (a) 300mV; (b) 600mV; (c) 900mV

In terms of propagation delay, the results suggest that ORDVC performs best. It obtains the maximum delay around the voltage close to the lower bound of CMR. MRDVC has the same trend as Aiello [36] that as the common-mode voltage goes up, the propagation delay increases first and then decreases. They both get a maximum delay at the voltages slightly below 1/2 *VDD*.

The relationship between the average power and common-mode voltage is obtained by providing a 5mV differential input and connecting a 5fF load. The clock frequency is set to 10MHz, 250MHz and 1GHz under the supply voltages of 300mV, 600mV and 900mV, respectively. It is reported in Fig. 4.5 that the structures based on the complementary comparators (MRDVC and Aiello [36]), always have maximum power consumption around 1/2 *VDD*. For the other two candidates adopting a single comparator (ORDVC and Seo [33]), they are little influenced by the common-mode voltage.



Fig. 4.5: Average power versus common-mode voltage under different supply voltages. (a) 300mV (b) 600mV (c) 900mV

Fig. 4.6 shows the impact of clock frequency on power consumption when common-mode voltage is biased to 1/2 VDD and a full-scale differential input is provided. Our proposed structures play a better performance on power than the other two candidates, which are dragged down by the complex circuit structure and shortcircuit current. Under the 300mV supply voltage, clock frequency does not greatly affect the power as leakage power consumption dominates. Under the 600mV and 900mV of VDD, every time the frequency is reduced by half, the power consumption of our proposed structures is also reduced by almost half, owing to the dynamic power dominating in the total power consumption. Since circuits do not always operate at the maximum clock frequency allowed by delay, our designs are attractive for power-saving in those applications with wide bandwidth. The trick to low power is that they control the dynamic power consumption caused by shortcircuit current in the decision phase.



Fig. 4.6: Average power versus clock frequency under different supply voltages. (a) 300mV; (b) 600mV; (c) 900mV

With Monte-Carlo simulation, the method in [38] is adopted to simulate the comparator offset. In our simulation setup, the maximum comparator offset variation that can be measured is around 1/3 *VDD*. Fig. 4.7 illustrates the impact of common-mode voltage on offset variation for the four candidates. As can be seen that all of the candidates achieve the minimum offset around 1/2 *VDD*. The best performers are MRDVC and Aiello [36], with a little difference between them. In the case of 300mV and 600mV supply voltage, ORDVC's offset variation is so large that it is out of the measurable range on the entire CMR. ORDVC also performed the worst under the 0.9V supply voltage.



Fig. 4.7: Comparator offset deviation versus common-mode voltage under different supply voltages. (a) 300mV; (b) 600mV; (c) 900mV

Fig. 4.8 shows the automatically generated layout through digital synthesis, demonstrating that the proposed structures are attractive for digital design. A comparison between the four dynamic voltage comparators is summarized in Tab. 4.1. It is noted that the data

derive from the simulations under the same simulation setup based on the 65nm CMOS process. Taken together, the simulation results suggest that MRDVC achieves a full rail-to-rail CMR and 4%~ 70% of power-saving(depending on clock frequency and supply voltage) compared to Seo [33] and Aiello [36], at the cost of less than 21% delay increment. ORDVC uses less than two-thirds of transistors, obtaining 45%~ 82% of power saving in contrast with the prior works. Despite more than twice the offset, ORDVC still has application scenarios, especially in stochastic flash ADC.



Fig. 4.8: Automatically generated layout based on 65nm CMOS process. (a)MRDVC with a latch; (b)ORDVC with a latch

Input offset [mV]	Power [nW]@{1.25 31.25 125} [MHz]	Power [nW]@{10 250 1000} [MHz]	CMR [mV]	Delay [ns]	V_{DD} [mV]	NO. transistors	Synthesizable		
49	14	26	150~300	12.35	300	72			2
41	240	1479	300~600	0.56	600		Yes	Seo[33]	
88	2099	14940	450~900	0.19	000				
29	8.3	12	0~90, 150~300	12.79	300	42		Aie	
34	217	612	0~600	0.60	600		Yes	llo[36]	
44	2727	6383	270~900	0.19	006				
>100	3.9	6.6	0~300	13.30	300	28		ORDVC	
>200	59	338	$180 \sim 480$	0.53	000		Yes		
177	489	3242	270~630	0.16	006				
28	5.5	10	0~300	15.49	300	46			
34	97	591	0~600	0.69	600		Yes	MRDVC	
49	805	5540	0~900	0.23	006			1C	

transistors (latches and buffers are included), propagation delay, CMR, power consumption and comparator offset variation Tab. 4.1: A comparison between the dynamic voltage comparators in regard to synthesizability, number of

4.3 Chapter Summary

In this chapter, we proposes a fully synthesizable rail-to-rail dynamic comparator, which can operate at a supply voltage down to 0.3V. Two variants are discussed and compared with other dynamic voltage comparators. Simulation demonstrates our proposed structures with optimum power efficiency under different supply voltages based on the 65nm CMOS process, which demonstrates that it can be competent in the SFADC with the requirements on low voltage and low power consumption.

CHAPTER 5

Linearity Enhancement Technique (LET)

5.1 Prior Works

The transfer function of the conventional SFADC has good linearity within the input range of $\pm l\sigma$, so generally SFADC' input range is between $\pm l\sigma$. On the one hand, the performance of the ADC usually improves as its linear input range increases. On the other hand, the non-linear transfer function will introduce harmonics into the system, thereby reducing the SNDR of the ADC. Therefore, in order to design an good ADC, we should improve the linearity of SFADC as much as possible and broaden the input range of ADC.

Some of linearity enhancement techniques has been proposed in the prior works. A technique [23] (shown as Fig. 5.1) has been presented, which reduces this nonlinearity by changing the overall transfer function by building a two-group SFADC. Setting the references of two comparator groups to have approximately $\pm l\sigma$ of comparator offsets allows higher linearity to be achieved. In this method, using the exact same comparators under the same conditions but merely dividing them into two groups with different references, an 8.5-dB improvement in SNDR can be obtained and there is no additional area overhead. Linearity is improved, however, the input range is still limited to $\pm l\sigma$.



Fig. 5.1: Two-group SFADC splitting the total number of comparators into two groups and applying an offset to each group, the shape of the transfer function can be controlled. For example, one group is given an offset of $+1\sigma$, and the other is -1σ .

The method above is further expanded in [39] (shown as Fig. 5.2). Comparators are equally divided into several comparator groups and reference voltages are equally spaced by a resistor ladder. Each group is referenced to a different reference voltage. Changing the mean (reference voltage) of comparator thresholds only shifts the transfer function along the input axis. Probability density functions (PDF) of comparator offset in each group is also shifted by the reference voltage. The overall output is obtained by summing the outputs of each group, achieving a flat top and a wide spread across the voltage range. While this method improves linearity, the input range is also greatly expanded.



Example of linearization for M = 8 case



Fig. 5.2: Multi-group SFADC

The prior work in [25] uses an additional hardware calibration circuit to perform piecewise linearization on the digital output, realizing a function similar to the inverse Gaussian function (shown as Fig. 5.3). The input range is extended to $[-3\sigma, 3\sigma]$. Although this algorithm is simple, it does not calibrate the higher-order terms in the transfer function and an additional hardware overhead for the calibration is needed.



Fig. 5.3: Inverse Gaussian CDF Method

In [40], to linearize the transfer function of the SFADC, a reference voltage is added to the comparator reference terminal, shown as Fig. 5.4. The added reference voltages follow random U-quadratic distribution. In this way, the trip point of the comparators will be spread to a wider range and shaped into uniform distribution. Let f(x), g(x) and h(x) be the PDF's of the comparators offset, U-quadratic reference signal and the resultant comparator trip points respectively. Thus $h(x) = g(x) \otimes f(x)$. In this method, the input range is extended to $[-3\sigma, 3\sigma]$. However, a large number of resistors with different values are used to generate the reference voltages following a quadratic distribution.



Fig. 5.4: U-quadratic distributed reference voltages method

5.2 Our proposed Linearity Enhancement Technique (LET)

In summary, it can be seen that the method of employing a multigroup architecture [39] can effectively extend the input range beyond $[-3\sigma, 3\sigma]$ and improve the linearity of the original transfer function to a certain extent. The above works [23,25,39,40] does not pay attention to the change of the comparator offset's standard deviation with the common mode voltage. When the reference voltages are different, although the distribution of the comparator offset also follows the Gaussian distribution, its shape has changed due to different standard deviation.

It will lead to a fluctuation at the top of total PDF if we simply combine these PDFs that are shifted by an equal space. Fig. 5.5 shows an example of SFADC consisting of multiply comparator groups where the number of groups is 11 and the reference voltage space is $l\sigma$. Even though we can change comparator offset through tuning the size of transistors in the comparator against the influence of common mode voltage, it is troublesome to customize the layout of comparators for the different groups. In addition, even if a constant comparator offset's standard deviation is guaranteed, there is still a limitation that the reference voltages should be set equally spaced.



Fig. 5.5: Total PDF when offset is not constant (Dark blue curve represents total PDF; the curves with other colors represent PDF of the groups).

5.2.1 Description of LET

To obtain a completely linear transfer function of SFADC, the total PDF of comparator offset should have a flat top. In this work, we propose a methodology that improves the linearity of SFADC through adjusting the weight of each group's PDF in the overall PDF. In the methodology, the comparator offset and reference voltage of the each group have more choices. Since SFADC represents probability with the yield of comparators outputting 'one', we can assign the number of comparators in each group to embody the weight. The value of the cumulative distribution function of the comparator offset also represents the probability that the input voltage is greater than the comparator offset. In SFADC, this probability is expressed by the proportion of the number of comparators that output 1 to the total. Thus

$$CDF_{Total} \cong P(V_{in} > V_{os}) = \frac{N_H}{N}$$
 (Eq. 5.1)

Where CDF_{total} , N_H , N are total PDF, the total number of comparators outputting '1' and the total number of comparators in SFADC, respectively.

Similarly,

$$CDF_i \cong \frac{N_{H,i}}{N_i}$$
 (Eq. 5.2)

Where CDF_i , $N_{H,i}$ and N_i are CDF of *i*-th group, the number of comparators outputting '1' in the *i*-th group and the number of comparators in *i*-th group, respectively.

Since

$$\sum N_i = N \quad \text{(Eq. 5.3)}$$

and

$$\sum N_{H,i} = N_H \quad \text{(Eq. 5.4)}$$

Combining (Eq. 5.1) to (Eq. 5.4), we can infer that

$$CDF_{Total} \cong \frac{N_H}{N} = \sum \frac{N_{H,i}}{N} = \sum \frac{N_i}{N} \cdot \frac{N_{H,i}}{N_i} \cong \sum \frac{N_i}{N} \cdot CDF_i$$
 (Eq.5.5)

Differentiate both sides of the formula (Eq. 5.5) we can get

$$PDF_{total} = \sum \frac{N_i}{N} \cdot PDF_i$$
 (Eq. 5.6)

It can be seen that the total PDF is the result of weighted summation of the PDFs in different comparator groups, and the weight is the proportion of the number of comparators in each group to the total number of comparators. Therefore, we can adjust the weight of the PDF in the overall PDF by controlling the proportion of each group in the overall comparators, so as to achieve a relatively flat PDF.

The conventional method is that, when the sigma of the offset hardly changes with the reference voltage, make each group have the same weight (comparator number) of PDF, and maintain a sigma distance between each PDF, so you can get a relatively flat at the top of the PDF. Therefore, the traditional method can be seen as a special case of our proposed scheme. When the standard deviation of the offset will be affected by the reference voltage, or the intervals between the reference voltages are not equidistant, our solution shows better applicability.

Hence, the issue of improving the linearity of SFADC is transformed into the linear combination of multiply PDFs, whose goal is to achieve a total PDF as flat as possible.

Assuming that we know the relationship of comparator offset variation versus reference voltage, we can get the probability density

64

function of input offset PDF_i in the *i*-th group under the reference voltage $V_{ref,i}$. The total PDF is derived as (w_i represents weight)

$$PDF_{total} = \sum w_i \cdot PDF_i$$
 (Eq. 5.7)
 $\sum w_i = 1$ and $w_i \ge 0$ (Eq. 5.8)

The fluctuation on the probability density of comparator offset straightway affects the linearity of transfer function. The best means to characterize the fluctuation is to analyze its AC power if we treat the PDF as a signal. The fluctuation can be evaluated by calculating the ratio of DC power to AC power (DAR). The greater the DAR is, the flatter the total PDF is. Given a fixed sampling interval (for example 0.1mV) and the input range(unit: mV) that we are interested in, we can calculate the DAR of total PDF in regard to weights w_i and PDF_i through the fast Fourier transform (FFT).

According to (Eq. 5.7), we can infer that

$$FFT(PDF_{total}) = \sum w_i \cdot FFT(PDF_i)$$
 (Eq. 5.9)

As we can see, when we change the weight, we do not have to reanalyze the FFT of total PDF. It can be obtained by weighting and summing the FFT of each PDF, which only needs to be analyzed once time. According to (Eq. 5.9), DC amplitude of total PDF (DC_{total}) and the AC amplitude of total PDF at the *j*-th spectral line $(AC_{total,j})$ can be calculated as

$$DC_{total} = \sum w_i \cdot DC_i$$
 (Eq. 5.10)

$$AC_{total,j} = \sum_{i=1}^{i=1} w_i \cdot AC_{i,j} \quad \text{(Eq. 5.11)}$$

where DC_i is the DC amplitude of *i*-th PDF, and $AC_{i,j}$ is the AC amplitude of *i*-th PDF at the *j*-th spectral, and *I* indicates the number of groups.

Therefore, according to the definition of DAR, we have

$$DAR = Power(DC_{total})/Power(AC_{total}) = DC_{total}^{2} / \sum_{j=2}^{j=N_{FFT}/2-1} AC_{total,j}^{2}$$
(Eq. 5.12)

where N_{FFT} indicates FFT points.

Substitute (Eq. 5.10) and (Eq. 5.11) into (Eq. 5.12), we can get

$$DAR = \left(\sum w_i \cdot DC_i\right)^2 / \sum_{j=2}^{j=N_{FFT}/2-1} \left(\sum_{i=1}^{i=I} w_i \cdot AC_{i,j}\right)^2$$
(Eq. 5.13)
= $F(w_1, w_2, ..., w_i)$

where DC_i and $AC_{i,j}$ have been obtained after the FFT of each PDF.

Thus, DAR is the function of weights $\{w_1, w_2, ..., w_i\}$. Through solving the optimum of multi-variable functions DAR under the constraints of (Eq. 5.8), we can determine the weight solution w_{opt} .

5.2.2 The reduction in the number of comparator groups

Initially we only impose the constraints (Eq. 5.8) on the weight solution and solve the global optimum DAR_{max} . Assuming that the weight solution corresponding to DAR_{max} is w_{opt} , we constrain those weight with a small value to be 0, and re-solve the global optimum value of DAR to see if it can be close to the initial DAR_{max} . If possible, we can say that these items with a weight of 0 are exclusion items, which has little effect on the flatness of the overall PDF across the range of interest. We can remove these terms from the overall PDF, which will help to reduce the number of comparator groups, thereby reducing the number of voltage references required.

5.2.3 The determination of the number of comparators

After obtaining the weights $\{w_1, w_2, ..., w_i\}$, assuming that the total number of comparators is N, then the number of comparators in the *i*-th group is allocated as

$$N_i = w_i / N$$
 (Eq. 5.14)

5.3 Effect of the proposed linearity enhancement technique

To verify the effect of the proposed linearity enhancement technique, we simulate the performance of the SFADC in behavior.

In order to determine the number of comparators in each group, we firstly sweep the the reference voltage and obtain the relationship between the comparator offset variation and the reference voltage through Monte Carlo simulation, as shown in Fig. 5.6. It can be seen that the standard deviation (1δ) of the comparator offset voltage

fluctuates around 20mV, so we choose 20mV as the reference voltage space. Theoretically, the SFADC based on MRDVC can reach the rail-to-rail input range. Nevertheless, under a certain number of MRDVCs, an increase in the input range will result in a decrease in the number of comparators per unit voltage range, thereby reducing resolution. In this case, we only focus on the input voltage range that we are interested in from 200mV to 400mV. The input range from 200mV to 400mV is divided into 10 equal parts through the resistor ladder and thus 11 reference voltages are created (namely 200mV:20mV:400mV).



Fig. 5.6: Standard deviation $(1 \ \delta)$ of comparator offset voltage versus reference voltage

V _{ref} [mV]	200	220	240	260	280	300	320	340	360	380	400
V _{os} [mV]	22.95	21.08	18.85	15.90	14.27	16.34	18.62	20.12	21.28	22.42	23.83
Proportion[%] (w/i LET)	20.8	0	11.3	6.5	6.7	7.7	8.6	6.7	11.2	0	20.5
Proportion[%] (w/o LET)	9.1	9.1	9.1	9.1	9.1	9.1	9.1	9.1	9.1	9.1	9.1

Tab. 5.1: The proportion of the number of comparators in each group

Through Monte Carlo simulation we can achieve the V_{os} (offset voltage's standard deviation) of MRDVC at these reference voltages. In addition, utilizing the method introduced in our proposed linearity enhancement technique, we can determine the proportion of the number of comparators in each group, shown in Tab. 5.1. As a control, the proportion of the number of comparators in each group when LET is not adopted is also revealed in Tab. 5.1.

Fig. 5.7 illustrates the total PDF without and with linearity enhancement technique respectively. It is obvious that the total PDF with LET has a flatter top than the other one. Fig. 5.8 and Fig. 5.9 reveal the distribution of offset randomly generated under the number of comparators N=1023 and N=4095, respectively. As can be seen that the offset's distribution with LET is more uniform within the voltage range from 200mV to 400mV, and the more the
number of comparators, the more obvious the effect of uniform distribution.

Fig. 5.10 shows the calculated ENOB in the case of different numbers of comparators when a full scale input is provided. Our method achieves a higher ENOB compared to the conventional method, with an improvements from 0.2 bits to 1.5 bits. As the total number of comparators increases, the improvement effect of ENOB compared to the conventional method is more obvious. Since the more comparators, the closer the actual comparator offset distribution in each group to the ideal Gaussian distribution, and the closer the total offset distribution to the ideal uniform distribution, thereby obtaining the improvement on ENOB compared to the conventional method.



Fig. 5.7: Total PDF without (Left) and with (Right) linearity enhancement technique



Fig. 5.8: Offset distribution without (Bottom) and with (Top) linearity enhancement technique when N=1023.



Fig. 5.9: Offset distribution without (Bottom) and with (Top) linearity enhancement technique when N=4095.



Fig. 5.10: ENOB versus the number of comparators

Fig. 5.11 shows the relationship between ENOB and the differential input voltage. It can be seen that when the differential voltage is less than 160mV, the ENOB of both methods increases with the increase of the input amplitude. The reason is that the increased input energy plays a major role. When the differential voltage amplitude exceeds 160mV, ENOB with our proposed method still maintains a positive correlation with the input voltage amplitude, and for the conventional method, ENOB begins to decrease with increasing voltage because the harmonic energy has exceeded the impact of the increased input energy on SNDR.



Fig. 5.11: ENOB versus input voltage magnitude (Vpp)

CHAPTER 6 Implementation and Simulation

We design an all-digital SFADC according to the configuration of Tab. 5.1 based on a 65nm CMOS process, as shown in Fig. 6.1, which is consist of 1023 MRDVCs, a 10-bits Wallace Tree Adder and a resistor ladder. All the cells constituting the comparator use the smallest cell in the standard cell library to obtain the largest comparator offset. All modules, except the resistor ladder, are described in verilog language.



Fig. 6.1: The structure of proposed SFADC, where k is equal to 11. Each group has different number of comparators inside, and there are 1023 comparators in total.

Fig. 6.2 shows the layout of the SFADC automatically generated by the digital synthesis process. The standard cells that make up the comparator and the ones that makes up the Wallace Tree Adder are placed together by the synthesis tool without any difference.



Fig. 6.2: Automatically generated Layout (250um*250um)

Under the 600mV supply voltage, with simulation we can measure the performance of the SFADC. Fig. 6.3 shows the output frequency spectrum of a sinusoidal input signal, whose frequency is 15.625MHz and amplitude is 100mV (full-scale), with a sample rate of 250MHz. The design achieves the SNDR of 35.3 dB and the SNR of 40.5 dB and the SFDR=38.6 dB, corresponding to the ENOB of 5.57 bits.



Fig. 6.3: Spectral power versus frequency(1024 points FFT)

Due to the difference in process and number of comparators, the power consumption among ADCs is unable to be compared straightly. A widely used FoM with respect to ENOB is introduced to estimate the overall performance of ADC, given as (Eq. 6.1). A comparison with other SFADCs in the reference paper [23], [25], [42], [43] and the flash ADC in [41] is shown in Tab. 6.1, which demonstrate the power effectiveness of our design. We can see our work plays the most excellent performance on FoM among reference SFADCs. Compared to the research work in [41], it significantly decreases the supply voltage.

FoM =
$$\frac{Power}{2^{\text{ENOB}} * f_s}$$
 (Eq. 6.1)

Tab. 6.1: Comparison of process, number of comparator, voltage, power in several kinds of SFADC

Reference	[23]	[25]	[41]	[42]	[43]	This work
Process	180nm	90nm	40nm	130nm	65nm	65nm
Ncomp	1152	2047	59	511	243	1023
ENOB	5.29	5.46	5.21	4.57	3.60	5.57
Voltage	900mV	700mV	1.1V	1.2V	600mV	600mV
Total	631uW	1.11mW	11mW	3.26mW	18.4uW	1.5mW
Power	@8MS/s	@21MS/s	@3GS/s	@100MS/s	@4MS/s	@250MS/s
FoM	2016fJ/con.	1201fJ/con.	100fJ/con.	1372fJ/con.	379fJ/con.	126fJ/con.

CHAPTER 7 Conclusion

This work presents a fully synthesizable stochastic flash A/D converter (SFADC), which can operate at the supply voltage of 0.6V with power consumption as low as 1.5mW at the clock frequency of 250MHz. By employing the all-digital comparator, the SFADC can be described with Verilog netlist and synthesized according to a standard digital design flow. Cross-coupled dynamic comparator structure saves the overall power due to remarkable control of addition. rail-to-rail dynamic power consumption. In the characteristic of comparator and the proposed linearity enhancement technique based on SFADC are proposed, allowing us to design a wide input-range stochastic flash ADC.

Our work revolves around the following. In the chapter 1, we introduce the trend of electronic device to low power consumption and miniaturization. Then we illustrate that the digital circuit enjoys the benefits of the upgrade of the process node more than the analog circuit. In view of the shortcomings of traditional flash ADC, we realize a digital stochastic flash ADC is a good idea to solve the defects.

In the chapter 2 we give a ADC's theoretical overview and introduce some of crucial specifications. In the chapter 3, we introduce the fundamental and properties of SFADC, and from the perspective of SNDR, theoretically establish a mathematical model of SNDR and the number of comparators required. For the cases in which SFADC contains a large number of comparator groups, Monte Carlo simulation is recommended to predict the number of comparators required.

In the chapter 4, in order to control the power consumption of the comparator which is used extensively in SFADC, we propose a fully synthesizable rail-to-rail dynamic comparator. We reveal its trick to control power consumption and the principle of realizing rail-to-rail input range. In addition, we discuss the two variants of it. Through the simulation, we compare them with other synthesizable dynamic voltage comparators.

In the chapter 5, we propose a linearity enhancement technique (LET) to improve the linearity and extend the input range of SFADC. We explain its principle and compare the performance of the SFADC with LET and the one without LET through behavior simulation. In the chapter 6, we design a fully synthesizable stochastic flash ADC based on the 65nm CMOS process and compare it to other prior works, showing the advantages in terms of supply voltage and power consumption.

The proposed SFADC is very suitable for the wearable system, such as health tracker, smart watch and communication system in

Internet of Things. Besides, it is also applicable in the biomedical field.

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